CSCI 210: Computer Architecture
Lecture 4: Introduction to MIPS

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Announcements

- Problem Set 1 due Friday 11:59 p.m.

Why you should learn (a little) assembly

- Learn what your computer is fundamentally capable of

- By learning about how high-level mechanisms are created in assembly, we learn what is fast, what is slow . . .

- Might use it for reverse engineering, embedded systems, compilers

CS History: Sophie Wilson

Developed the ARM Instruction Set Architecture

The MIPS Instruction Set

- Used as the example throughout the book
- Stanford MIPS commercialized by MIPS Technologies (owned by John L. Hennessy, who wrote your book.)
- Used in Embedded Systems
  - Applications in consumer electronics, network/storage equipment, cameras, printers, …
- Typical of many modern ISAs
  - Most similar to ARM, RISC-V

Three Types of Instruction

- Arithmetic and logical (R)
  - Operates on data entirely in registers
- Immediate (I)
  - One of the operands is encoded directly in the instruction
- Jump (J)
  - Changes the pc to a new location

Arithmetic and Logical Operations

- Add and subtract, three operands
  - Two sources and one destination

```
add a, b, c  # a = b + c
sub a, b, c  # a = b - c
and a, b, c  # a = b & c (bit-wise AND)
```

- All arithmetic and logical operations have this form

Convert to MIPS: f = (g + h) - (i + j);

A.

```
add   f, g, h
sub   f, i, j
```

B.

```
add   t0, g, h
add   t1, i, j
sub   f, t0, t1
```

C.

```
sub   f, (add g,h), (add i,j)
```

D. More than one of these is correct

```
add a, b, c  # a = b + c
sub a, b, c  # a = b - c
```

Register Operands

- Arithmetic instructions use register operands

- MIPS has 32 32-bit general purpose registers
  - Numbered 0 to 31
  - 32-bit data called a “word”

- ARM has 37 32-bit general purpose registers
- X86-64 has 16 general purpose registers, around 40 named registers used by the processor
  - Can be used as 8, 16, 32, or 64 bit registers

Aside: MIPS Register Convention

Name	Register Number	Usage
$zero	0	constant 0 (hardware)
$at	1	reserved for assembler
$v0–$v1	2–3	returned values
$a0–$a3	4–7	arguments
$t0–$t7	8–15	temporaries
$s0–$s7	16–23	saved values
$t8–$t9	24–25	temporaries
$gp	28	global pointer
$sp	29	stack pointer
$fp	30	frame pointer
$ra	31	return addr (hardware)

Register Operand Example

- C code:

```
f = (g + h) - (i + j);
```

  – f, g, h, and j in registers `$s0`, `$s1`, `$s2`, `$s3`, and `$s4`

- Compiled MIPS code:

```
add  $t0, $s1, $s2
add  $t1, $s3, $s4
sub  $s0, $t0, $t1
```

Some R-type instructions

```
• add  dest, src1, src2
• sub  dest, src1, src2
• mul  dest, src1, src2   # Pseudoinstruction!
• div  dest, src1, src2   # Pseudoinstruction!
• move dest, src          # add dest, $zero, src
• and  dest, src1, src2
• or   dest, src1, src2
• nor  dest, src1, src2
• xor  dest, src1, src2
```

Assume registers initially have the following values

$a0	$a1	$t0	$t1	$v0
2	100	5	6	7

What values do they have after running this code?

```
move $t0, $a0
add  $t1, $a0, $a0
add  $t1, $t1, $t1
sub  $t0, $t1, $t0
add  $v0, $t0, $a1
```

	$a0	$a1	$t0	$t1	$v0
A	2	100	5	6	7
B	2	100	6	8	106
C	5	-10	-17	22	7
D	5	100	15	20	115
E	None of the above				

Questions about Arithmetic Operations?

Memory Instructions

- `lw   $t0, 0($t1)`
  - `$t0 = Mem[$t1+0]`
  - Loads 4 bytes from `$t1, $t1+1, $t1+2, and $t1+3`
- `sw   $t0, 4($t1)`
  - `Mem[$t1+4] = $t0`
  - Stores 4 bytes at `$t1+4,  $t1+5,  $t1+6, and $t1+7`
- These instructions are the cornerstones of our being able to move data to and from memory

Load instructions

- **`lw`** — Loads 4 bytes of memory into a register
  - `lw    $t0, 8($t4)`
- **`lh`** — Loads 2 bytes of memory into a register
  - `lh    $t2, 6($t1)`
- **`lb`** — Loads 1 byte of memory into a register
  - `lb    $t3, 3($t0)`

- `lw` and `lb` are more common than `lh`

Store instructions

- **sw** — Stores 4 bytes from a register into memory
  - `sw    $t0, 8($t4)`
- **sh** — Stores 2 bytes from a register into memory
  - `sh    $t2, 6($t1)`
- **sb** — Stores 1 byte from a register into memory
  - `sb    $t3, 3($t0)`

- `sw` and `sb` are more common than sh

Accessing the Operands

There are typically two locations for nonconstant operands – registers (internal storage e.g., $t0 or $a0) and memory. In each column we have which—reg or mem—is better. Which row is correct?

	Faster access	Smaller number to specify a reg/mem location	More locations
A	Mem	Mem	Reg
B	Mem	Reg	Mem
C	Reg	Mem	Reg
D	Reg	Reg	Mem
E	None of the above		

Load-store architectures

can do:

```
load r3, M(address)
add r1 = r2 + r3
```

$\Rightarrow$ forces heavy dependence on registers, which is exactly what you want in today's CPUs

can't do

```
add r1 = r2 + M(address)
```

- more instructions

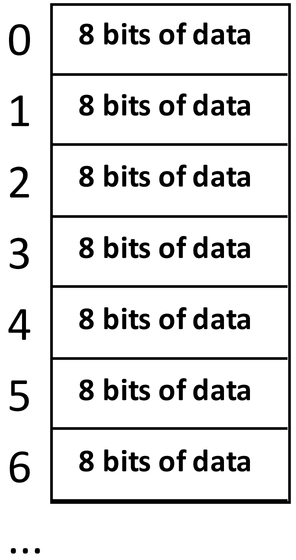
+ fast implementation

Memory

- Main memory used for composite data
  - Arrays, structures, dynamic data
- Memory is byte addressed
  - Each address identifies an 8-bit byte
- Words are **aligned** in memory
  - Address of a word must be a multiple of 4
  - A word whose address is not a multiple of 4 is **misaligned**
  - Misaligned memory accesses cause a hardware exception in MIPS

Memory Organization

- Viewed as a large, single-dimension array
- A memory address is an index into this array
- “Byte Addressing" means that the index points to a byte of memory.



Memory Organization

- Bytes are nice, but most data items use larger "words"
- For MIPS, a word is 32 bits or 4 bytes.

0	**32 bits of data**
4	**32 bits of data**
8	**32 bits of data**
12	**32 bits of data**

Registers hold 32 bits of data

- 2^{32} bytes with byte addresses from 0 to 2^{32} - 1
- 2^{30} words with byte addresses 0, 4, 8, ... 2^{32} - 4

If you have a pointer to address 1000 and you increment it by one to 1001. What does the new pointer point to, relative to the original pointer?

A) The next word in memory
B) The next byte in memory
C) Either the next word or byte – depends on if you use that address for a load byte or load word
D) Pointers are a high level construct – they don't make sense pointing to raw memory addresses.
E) None of the above.

If a 4-byte word is in memory at address 4203084, what is the address of the next word in memory?

A) 4203085
B) 4203088
C) 14203084
D) It depends on the value of the words in memory
E) Since a word is 4 bytes, it’s not possible to have one at address 4203084

Getting the address of data in the first place

- Three main locations for data in a program
  - **Global variables (these live in the data segment)**
  - Local variables (function call stack)
  - Dynamically allocated memory (memory allocated at runtime)

Global variables

- Global variables live in the data segment
- We use assembler directives to
  1. Switch to the data segment
  2. Allocate space for the globals
  3. Switch back to the text (code) segment

```
.data # Switches to data segment
nums: # Label we can use to get the address of the following array
.word 37, -42, 806 # allocates space for 3 words
.text # Switches back to the text segment
```

Load address pseudo instruction

Sets a register to the address pointed to by the symbolic label

```
.data
nums: .word 37, -42, 806

.text
main:
  la   $s0, nums
  lw   $t0, 0($s0)
  lw   $t1, 4($s0)
  lw   $t2, 8($s0)
```

Reading

- Next lecture: Assembly
  - 2.3

- Problem Set 1: Due Friday at 11:59 p.m.